

# CS 24

## Introduction to Computing Systems



# Outline

1 Compilation and JVM

2 Memory

3 Integers

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Hardware

```
01010101  
11011011  
01110111
```

Machine Code

```
movl %eax, -12(%rbp)  
movl %esi, %eax  
addq $16, %rsp
```

x86-64 Assembly

```
char *x = malloc(sizeof(char));  
*x = 'a';  
printf("%c\n", x);  
free(x);
```

C Program



Java Virtual Machine

```
64 05 6c
```

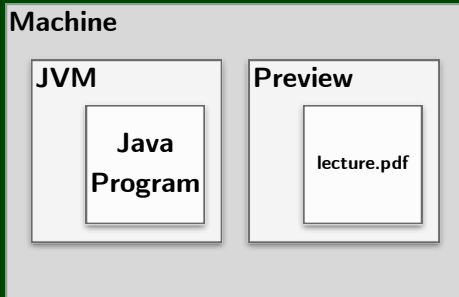
Java Bytecode

```
isub  
iconst_2  
idiv
```

Java Bytecode  
Instructions

```
public class Hello {  
    public static void main(String[] args) {  
        System.out.println("Hello");  
    }  
}
```

Java Program



## Overview

In this project, you will implement all the integer JVM instructions. Your JVM will be able to run **real** compiled class files.

## Learning Outcomes

- I can distinguish between how Java and C execute on a computer.
- I can identify the different levels of expressiveness between assembly/bytecode and statements in a high-level programming language.
- I can describe how code can be viewed as a type of data.
- I can write a virtual machine.

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## Memory, Addresses, and Pointers

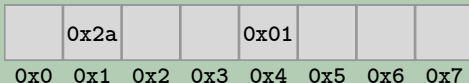
- **Memory** is (essentially) a large array of bytes.
- An **address** is an index into that array.
- A **pointer** is a variable that stores an address.

```
1 char *p = malloc(sizeof(char));
2 *p = 42;
3 printf("p = %p\n", p);
4 printf("*p = %p\n", *p);
5 printf("&p = %p\n", &p);
```

### OUTPUT

```
>> p = 0x01
>> *p = 0x2a
>> &p = 0x04
```

## A Picture of Memory

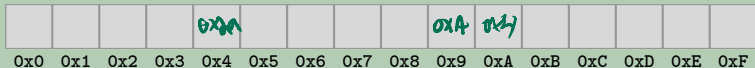


```
1 char **p = malloc(sizeof(char *));  
2 *p = malloc(sizeof(char));  
3 **p = 42;  
4 printf("p = %p\n", p);  
5 printf("*p = %p\n", *p);  
6 printf("**p = %p\n", **p);  
7 printf("&p = %p\n", &p);  
8 printf("&*p = %p\n", &*p);  
9 printf("&**p = %p\n", &**p);
```

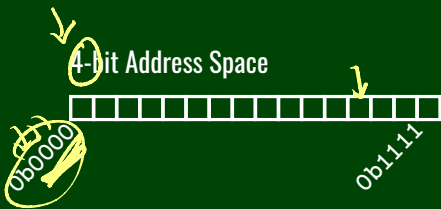
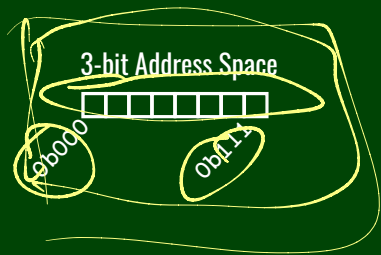
OUTPUT

```
>> p = 0x0a  
>> *p = 0x04  
>> **p = 0x2a  
>> &p = 0x09  
>> &*p = 0x0a  
>> &**p = 0x0a
```

## A Picture of Memory



□ = 1 byte



3-bit Address Space

4-bit Address Space

5-bit Address Space

6-bit Address Space

7-bit Address Space

8-bit Address Space

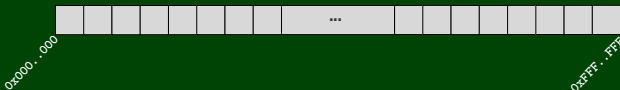
## Poll

How many bits are necessary to represent an address in a tiny computer with only 8 addressable bytes?

- a 3
- b 4
- c 8
- d 64
- e ???

The **word size** of a machine is the size of its **registers** and **addresses**.

labradoodle (and most other machines) have a 64-bit word size. This gives us 18 EB (exabytes) of addressable memory.

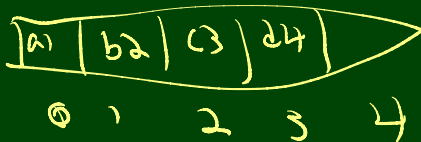


To reference a word, we use the address **of the first byte**. Thus, to move to the next word, we add **eight** (64-bit register = 8 bytes).

So, how are the bytes **within** a multi-byte word ordered in memory?

OUTPUT

```
>> x = 0xa1b2c3d4  
>> &x = 0x100
```



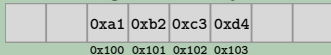
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OUTPUT

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>> x = 0xa1b2c3d4  
>> &x = 0x100
```

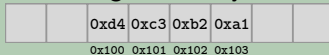
Big Endian (Internet, JVM)

**Most** Significant Byte First



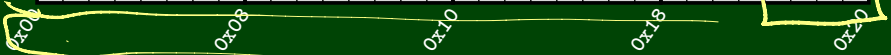
Little Endian (x86, ARM <sup>(most OSes)</sup>)

**Least** Significant Byte First



```
1 uint8_t *p1 = 16;  
2 uint32_t *p2 = 0x1C;
```

FFCA0110DFEBCAFE2983287323622E8EFFFFFFDDEAFACE7EE8D9A64E8C000000



What are the values of `*p1` and `*p2` (in decimal) on a **little endian** machine?

`uint8_t *p = ~~~~~`  
`*p`