

CS
24

Introduction to Computing Systems

x86-64 Basics

```
mov %edi, %eax
xor %ecx, %ecx
test %edi, %edi
setne %cl
mov $0xffffffff, %eax
cmovns %ecx, %eax
retq
```

- In 2021, almost all assembly is written by **compilers** or in the **operating system**, but who writes those?
- In 2021, the high-level language model occasionally breaks down and you have to read the assembly to understand the machine's behavior!
- In 2021, it's important to understand the types of optimizations a compiler is capable of making—and those it isn't!
- In 2021, software is generally distributed in binary form; if you want to **reverse engineer** or **security audit** software, it's going to be assembly!

Compiling a Program

```
blank@labradoodle:~$ cat identity.c
int identity(int x) {
    return x;
}
```

```
blank@labradoodle:~$ clang -S identity.c
blank@labradoodle:~$ cat identity.s
identity:
    movl %edi, %eax
    retq
```

```
blank@labradoodle:~$ cat identity.c
int identity(int x) {
    return x;
}
```

```
blank@labradoodle:~$ clang -c identity.c
blank@labradoodle:~$ objdump -d identity.o
```

```
simple.o:      file format elf64-x86-64
```

Disassembly of section .text:

```
0000000000000000 <identity>:
 0: 89 f8          mov    %edi,%eax
 2: c3             retq
```

- **Immediates:** Constant integer data
 - Examples: \$0x400, \$-533
 - Like C literal, but prefixed with '\$'
 - Encoded with 1, 2, 4, or 8 bytes depending on the instruction
- **Registers:** behave like “global variables”, but hardwired in the processor
 - Examples: %eax, %edi
 - Some of them are reserved for special uses or have special meanings
- **Memory:** Consecutive bytes of memory

What are these %eax things?

Registers are locations in the CPU that store a small amount of data, which can be accessed very quickly (once every clock cycle). They have names (e.g., %rsi)—not addresses. They are a precious commodity in all architectures, but especially x86-64.

(return) %eax

%ebx

(arg 4) %ecx

(arg 3) %edx

(arg 1) %edi

(arg 2) %esi

%esp

%ebp

There are three major **classes** of things assembly instructions do:

- 1 Transfer data between memory and registers
 - Load data from memory into register: `%reg = mem[addr]`
 - Store register data into memory: `mem[addr] = %reg`
- 2 Perform arithmetic operation on register or memory data
 - e.g., `%eax += %ebx`
 - e.g., `%eax += mem[addr]`
- 3 Re-direct control flow (jumps and gotos)

Understanding Identity (`mov`)

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C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

Understanding Identity (`mov`)

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C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
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```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

Pseudocode Translation (so far)

```
1 identity:  
2     %eax = %edi  
3     retq
```

Understanding Identity (`retq` and return value)

C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

↔

C Pseudocode

`return %eax`

Understanding Identity (`retq` and return value)

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C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

C Pseudocode

`return %eax`

Pseudocode Translation (so far)

```
1 identity:  
2     %eax = %edi  
3     return %eax
```

Understanding Identity (arguments)

C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

↔

C Pseudocode

`return %eax`

Understanding Identity (arguments)

C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

↔

C Pseudocode

`return %eax`

Pseudocode Translation (so far)

```
1 identity(%edi):  
2     %eax = %edi  
3     return %eax
```

Understanding Identity (simplifying)

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C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
400482: c3             retq
```

x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

↔

C Pseudocode

`return %eax`

Understanding Identity (simplifying)

10

C Code

```
1 int identity(int x) {  
2     return x;  
3 }
```

x86-64 Disassembly

```
0000000000400480 <identity>:  
400480: 89 f8          mov    %edi,%eax  
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x86-64 Instruction: `mov`

x86-64

`mov %src, %dst`

↔

C Pseudocode

`%dst = %src`

x86-64 Instruction: `retq`

x86-64

`retq`

↔

C Pseudocode

`return %eax`

Pseudocode Translation (so far)

```
1 identity(%edi):  
2     return %edi
```

System V AMD64 ABI

- The value in %eax is automatically returned by `retq`.
- The first argument to a function is stored in %edi.

Things to Notice About x86-64

- There are no types!!!
- The **conventions** define what the compiler is allowed to do.

Understanding Arithmetic (arithmetic instructions)

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```
0000000000000000 <arith>:  
0: 83 c7 01          add    $0x1,%edi  
3: 0f af ff          imul   %edi,%edi  
6: 89 f8            mov    %edi,%eax  
8: c3                retq
```

```
0000000000000000 <arith>:  
0: 83 c7 01          add    $0x1,%edi  
3: 0f af ff          imul   %edi,%edi  
6: 89 f8            mov    %edi,%eax  
8: c3                retq
```

Pseudocode Translation (so far)

```
1 arith(%edi):  
2     add $0x1,%edi  
3     imul %edi,%edi  
4     %eax = %edi  
5     return %eax
```

```
0000000000000000 <arith>:  
0: 83 c7 01          add    $0x1,%edi  
3: 0f af ff          imul   %edi,%edi  
6: 89 f8            mov    %edi,%eax  
8: c3                retq
```

Pseudocode Translation (so far)

```
1 arith(%edi):  
2     add $0x1,%edi  
3     imul %edi,%edi  
4     %eax = %edi  
5     return %eax
```

x86-64 Instruction: (Some) Arithmetic Operations

x86-64

add %src, %dst	↔
sub %src, %dst	↔
imul %src, %dst	↔

C Pseudocode

%dst += %src
%dst -= %src
%dst *= %src

```
0000000000000000 <arith>:  
0: 83 c7 01          add    $0x1,%edi  
3: 0f af ff          imul   %edi,%edi  
6: 89 f8            mov    %edi,%eax  
8: c3                retq
```

x86-64 Instruction: (Some) Arithmetic Operations

x86-64

add %src, %dst	↔
sub %src, %dst	↔
imul %src, %dst	↔

C Pseudocode

%dst += %src
%dst -= %src
%dst *= %src

Pseudocode Translation (so far)

```
1 arith(%edi):  
2     %edi += $0x1  
3     %edi *= %edi  
4     %eax = %edi  
5     return %eax
```

```
0000000000000000 <arith>:  
0: 83 c7 01          add    $0x1,%edi  
3: 0f af ff          imul   %edi,%edi  
6: 89 f8            mov    %edi,%eax  
8: c3                retq
```

x86-64 Instruction: (Some) Arithmetic Operations

x86-64

add %src, %dst	↔
sub %src, %dst	↔
imul %src, %dst	↔

C Pseudocode

%dst += %src
%dst -= %src
%dst *= %src

Pseudocode Translation (so far)

```
1 arith(%edi):  
2     return (%edi + $0x1) * (%edi + $0x1)
```

Things to Notice

- There are no types!!!
- The **ABI** defines what the compiler should do.

(return)

%eax

(arg 4)

%ebx

(arg 3)

%ecx

(arg 1)

%edx

(arg 2)

%edi

%esi

%esp

%ebp

Another Mystery Function (bit operations)

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```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1        setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1        cmovns %ecx,%eax  
f: c3             retq
```

(return) %eax

%ebx

%ecx

(arg 4)

%edx

(arg 3)

%edi

(arg 1)

%esi

(arg 2)

%esp

%ebp

Pseudocode Translation (so far)

Another Mystery Function (bit operations)

16

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1        setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1        cmovns %ecx,%eax  
f: c3             retq
```

(return) %eax
 %ebx
 %ecx
 %edx
 (arg 1) %edi
 (arg 2) %esi
 %esp
 %ebp

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ...
```

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1       setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1       cmovns %ecx,%eax  
f: c3             retq
```

x86-64 Instruction: `test`

The processor has a special register that contains “flags” which `test` sets.

`test %r1, %r2`

- ZF set to result of $(%r1 \& %r2) == 0$
- SF set to result of $(%r1 \& %r2) < 0$

Another Mystery Function (`test`)

17

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1        setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1        cmovns %ecx,%eax  
f: c3             retq
```

(return) %eax

%ebx

(arg 4) %ecx

(arg 3) %edx

(arg 1) %edi

(arg 2) %esi

%esp

%ebp

CF

ZF

SF

OF

x86-64 Instruction: `test`

The processor has a special register that contains “flags” which `test` sets.

`test %r1, %r2`

- ZF set to result of $(%r1 \& %r2) == 0$
- SF set to result of $(%r1 \& %r2) < 0$

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ZF = %edi == 0  
4     SF = %edi < 0  
5     ...
```

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1        setne  %cl  
7: b8 ff ff ff ff  mov    $0xffffffff,%eax  
c: 0f 49 c1        cmovns %ecx,%eax  
f: c3              retq
```

x86-64 Instruction: set??

x86-64

sete %r

↔

%r = ZF

setne %r

↔

%r = ~ZF

C Pseudocode

sets %r

↔

%r = SF

setns %r

↔

%r = ~SF

Another Mystery Function (`set??`)

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```
0000000000000000 <mystery>:  
0: 31 c9           xor    %ecx,%ecx  
2: 85 ff           test   %edi,%edi  
4: 0f 95 c1         setne  %cl  
7: b8 ff ff ff ff  mov    $0xffffffff,%eax  
c: 0f 49 c1         cmovns %ecx,%eax  
f: c3              retq
```

(return) %eax

%ebx

(arg 4) %ecx

(arg 3) %edx

(arg 1) %edi

(arg 2) %esi

%esp

%ebp

CF

ZF

SF

OF

x86-64 Instruction: `set??`

x86-64

<code>sete %r</code>	↔	$\%r = \text{ZF}$
<code>setne %r</code>	↔	$\%r = \sim\text{ZF}$
<code>sets %r</code>	↔	$\%r = \text{SF}$
<code>setns %r</code>	↔	$\%r = \sim\text{SF}$

C Pseudocode

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ZF = %edi == 0  
4     SF = %edi < 0  
5     %ecx = ~ZF  
6     ...
```

Another Mystery Function (`mov`)

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```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1       setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1       cmovns %ecx,%eax  
f: c3             retq
```

Another Mystery Function (`mov`)

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```
0000000000000000 <mystery>:  
0: 31 c9           xor    %ecx,%ecx  
2: 85 ff           test   %edi,%edi  
4: 0f 95 c1         setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1         cmovns %ecx,%eax  
f: c3              retq
```

(return) %eax

%ebx

(arg 4) %ecx

(arg 3) %edx

(arg 1) %edi

(arg 2) %esi

%esp

%ebp

CF

ZF

SF

OF

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ZF = %edi == 0  
4     SF = %edi < 0  
5     %ecx = ~ZF  
6     %eax = -1  
7     ...
```

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1       setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1       cmovns %ecx,%eax  
f: c3             retq
```

x86-64 Instruction: `cmov??`

x86-64

<code>cmove</code> %src, %dst	\leftrightarrow	%dst = %src if ZF
<code>cmovne</code> %src, %dst	\leftrightarrow	%dst = %src if ~ZF
<code>cmovs</code> %src, %dst	\leftrightarrow	%dst = %src if SF
<code>cmovns</code> %src, %dst	\leftrightarrow	%dst = %src if ~SF

C Pseudocode

Another Mystery Function (`cmov??`)

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```
0000000000000000 <mystery>:  
0: 31 c9           xor    %ecx,%ecx  
2: 85 ff           test   %edi,%edi  
4: 0f 95 c1         setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1         cmovns %ecx,%eax  
f: c3              retq
```

(return) %eax

(arg 4) %ebx

(arg 3) %ecx

(arg 1) %edx

(arg 2) %edi

(arg 1) %esi

(arg 2) %esp

(arg 3) %ebp

x86-64 Instruction: `cmov??`

x86-64

<code>cmove</code> %src, %dst	\leftrightarrow	%dst = %src if ZF
<code>cmovne</code> %src, %dst	\leftrightarrow	%dst = %src if ~ZF
<code>cmovs</code> %src, %dst	\leftrightarrow	%dst = %src if SF
<code>cmovns</code> %src, %dst	\leftrightarrow	%dst = %src if ~SF

C Pseudocode

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ZF = %edi == 0  
4     SF = %edi < 0  
5     %ecx = ~ZF  
6     %eax = -1  
7     if (~SF) { %eax = %ecx }  
8     ...
```

CF ZF SF OF

Another Mystery Function (`retq`)

21

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1       setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1       cmovns %ecx,%eax  
f: c3             retq
```

Another Mystery Function (`retq`)

21

```
0000000000000000 <mystery>:  
0: 31 c9          xor    %ecx,%ecx  
2: 85 ff          test   %edi,%edi  
4: 0f 95 c1       setne  %cl  
7: b8 ff ff ff ff mov    $0xffffffff,%eax  
c: 0f 49 c1       cmovns %ecx,%eax  
f: c3             retq
```

(return)	%eax
	%ebx
(arg 4)	%ecx
(arg 3)	%edx
(arg 1)	%edi
(arg 2)	%esi
	%esp
	%ebp
	CF ZF SF OF

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %ecx = 0  
3     ZF = %edi == 0  
4     SF = %edi < 0  
5     %ecx = ~ZF  
6     %eax = -1  
7     if (~SF) { %eax = %ecx }  
8     return %eax
```

Another Mystery Function (simplifying)

22

```
0000000000000000 <mystery>:  
0: 31 c9           xor    %ecx,%ecx  
2: 85 ff           test   %edi,%edi  
4: 0f 95 c1         setne  %cl  
7: b8 ff ff ff ff  mov    $0xffffffff,%eax  
c: 0f 49 c1         cmovns %ecx,%eax  
f: c3              retq
```

(return) %eax

%ebx

(arg 4) %ecx

(arg 3) %edx

(arg 1) %edi

(arg 2) %esi

%esp

%ebp

CF

ZF

SF

OF

Pseudocode Translation (so far)

```
1 mystery(%edi):  
2     %eax = -1  
3     if (%edi >= 0) { %eax = %edi != 0 }  
4     return %eax
```